



# Enhancing Key Management Process in Wireless Sensor Network

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*Abstract—This paper shows the certificateless key management protocol for secure communication in wireless sensor network. This paper proposes a certificateless-effective key management (CL-EKM) protocol for secure communication in dynamic WSNs characterized by node mobility. The CL-EKM supports efficient key updates when a node leaves or joins a cluster and ensures forward and backward key secrecy. Many cluster-based wireless sensor network routing protocols have been proposed. However, most of them take little consideration on communication protection, which is important to ensure the network security. The CL-EKM supports efficient key updates when a node leaves or joins a cluster and ensures forward and backward key secrecy. wireless sensor systems (WSNs) have been conveyed for a wide assortment of utilizations, including military detecting and following, tolerant status observing, activity stream checking, where tactile gadgets regularly move between distinctive areas. Securing information and interchanges requires suitable encryption key conventions*

*Keywords— Wireless sensor networks, certificateless public key cryptography, elliptic curve cryptography (ECC), cluster.*

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## I. INTRODUCTION

Dynamic wireless sensor networks (WSNs), which enable mobility of sensor nodes, facilitate wider network coverage and more accurate service than static WSNs. Therefore, dynamic WSNs are being rapidly adopted in monitoring applications, such as target tracking in battlefield surveillance, healthcare systems, traffic flow and vehicle status monitoring, dairy cattle health monitoring [9]. However, sensor devices are vulnerable to malicious attacks such as impersonation, interception, capture or physical destruction, due to their unattended operative environments and lapses of connectivity in wireless communication. Thus, security is one of the most important issues in many critical dynamic WSN applications. Dynamic WSNs thus need to address key security requirements, such as node authentication, data confidentiality and integrity, whenever and wherever the nodes move.

This paper present a certificateless effective key management (CL-EKM) scheme for dynamic WSNs. In certificateless public key cryptography (CL-PKC) [12], the user's full private key is a combination of a partial private key generated by a key generation center (KGC) and the user's own secret value. The special organization of the full

private/public key pair removes the need for certificates and also resolves the key escrow problem by removing the responsibility for the user's full private key. We also take the benefit of ECC keys defined on an additive group with a 160-bit length as secure as the RSA keys with 1024-bit length.

## II. LITRATURE SERVEY

We propose the first certificateless effective key management protocol for secure communication in dynamic WSNs. Certificateless effective key management supports efficient communication for key updates and management when a node leaves or joins a cluster and hence ensures forward and backward key secrecy [1]. Wireless sensor networks come with huge application domain but on the other hand require the same level of security.

The paper discusses various authentication techniques available in wireless sensor network and analyzes them. Some techniques are very helpful but come with some disadvantages. The effort is also done to point out these difficulties. Authentication is one of the best security solutions which protects whole sensor network.[2] ]. We have identified wireless sensor network applications, classified sensor networks into different classes and identified security attacks that can take place in each class of sensor networks.[7]

## III.CERTIFICATELESS EFFECTIVE KEY MANAGEMENT SCHEME

### Types of Keys

- **Certificateless Public/Private Key:** Before a node is deployed, the KGC at the BS generates a unique certificateless private/public key pair and installs the keys in the node. This key pair is used to generate a mutually authenticated pairwise key.
- **Individual Node Key:** Each node shares a unique individual key with BS. For example, a L-sensor can use the individual key to encrypt an alert message sent to the BS, or if it fails to communicate with the H-sensor. An H-sensor can use its individual key to encrypt the message corresponding to changes in the cluster. The BS can also use this key to encrypt any sensitive data, such as compromised node information or commands. Before a node is deployed, the BS assigns the node the individual key.
- **Pairwise Key:** Each node shares a different pairwise key with each of its neighboring nodes for secure communications and authentication of these nodes. For example, in order to join a cluster, a L-sensor should share a pairwise key with the H-sensor. Then, the H-sensor can securely encrypt and distribute its cluster key to the L-sensor by using the pairwise key. In an aggregation supportive WSN, the L-sensor can use its pairwise key to securely transmit the sensed data to the H-sensor. Each node can dynamically establish the pairwise key between itself and another node using their respective certificateless public/private key pairs.
- **Cluster Key:** All nodes in a cluster share a key, named as cluster key. The cluster key is mainly used for securing broadcast messages in a cluster, e.g., sensitive commands or the change of member status in a cluster. Only the cluster head can update the cluster key when a L-sensor leaves or joins the cluster.

#### IV. WORKING OF PROPOSE KEY MANAGEMENT SYSTEM

##### A. system setup:

1) Generation of System Parameters : The KGC at the BS runs the following steps by taking a security parameter  $k \in \mathbb{Z}^+$  as the input, and returns a list of system parameter  $\_ = \{Fq, E/Fq, Gq, P, Ppub = xP, h0, h1, h2, h3\}$  and  $x$ .

2) Node Registration : The BS assigns a unique identifier, denoted by  $Li$ , to each L-sensor  $nLi$  and a unique identifier, denoted by  $Hj$ , to each H-sensor  $nHj$ , where  $1 \leq i \leq N1$ ,  $1 \leq j \leq N2$ ,  $N = N1 + N2$ . Here we describe the certificateless public/private key and individual node key operations for  $Li$ , the same mechanisms apply for H-sensors. During initialization, each node  $nLi$  chooses a secret value  $xLi \in \mathbb{R} \mathbb{Z} * q$  and computes  $PLi = xLi P$ . Then, the BS requests the KGC for partial private/public keys of  $nLi$  with the input parameters  $Li$  and  $PLi$ . The KGC chooses  $rLi \in \mathbb{R} \mathbb{Z} * q$  and then computes a pair of partial public/private key  $(RLi, dLi)$  as below:

$$RLi = rLi P$$

$$dLi = rLi + x \cdot h0(Li, RLi, PLi) \text{ mod } q.$$

##### B. Pairwise Key Generation:

1) Pairwise Master Key Establishment: In this paragraph, we describe the protocol for establishing a pairwise master key between any two nodes  $nA$  and  $nB$  with unique IDs  $A$  and  $B$ , respectively. We utilize the CL-HSC scheme [13] as a building block. When  $nA$  receives an advertisement message from  $nB$ , it executes the following encapsulation process to generate a long-term pairwise master key  $KAB$  and the encapsulated key information,  $\phi A = (UA, WA)$ .

2) Pairwise Encryption Key Establishment : Once  $nA$  and  $nB$  set the pairwise master key  $KAB$ , they generate an HMAC of  $KAB$  and a nonce  $r \in \mathbb{R} \mathbb{Z} * q$ . The HMAC is then validated by both  $nA$  and  $nB$ . If the validation is successful, the HMAC value is established as the short-term pairwise encryption key  $kAB$ .

##### C. Cluster Formation:

1) Node Discovery and Authentication: For node discovery,  $nHj$  broadcasts an advertisement message containing  $Hj$  and  $pkHj$ . Once  $nLi$  within  $Hj$ 's radio range receives the advertisement, it checks  $Hj$  and  $pkHj$ , and initiates the Pairwise Key Generation procedure. Note that  $nLi$  may receive multiple advertisement messages if it is within the range of more than one H-sensor. However,  $nLi$  must choose one H-sensor, may be by prioritizing over the proximity and signal strength. Additionally,  $nLi$  can record other H-sensor advertisements as backup cluster heads in the event that the primary cluster head is disabled. If  $nLi$  selects multiple cluster heads and sends a response to all of them, it is considered as a compromised node.  $nLi$  and  $nHj$  perform the Pairwise Key Generation procedure to obtain a pairwise master key,  $KLi Hj$  and a pairwise encryption key,  $kLi Hj$ .

2) Cluster Key Generation:  $nHj$  chooses  $xj \in \mathbb{R} \mathbb{Z} * q$  to generate a cluster key  $GKj$  as follows  $GKj = \text{HMAC}(xj, Hj)$ . Then,  $nHj$  computes  $C2 = EkLi Hj (GKj, Hj, Li)$  to distribute the  $GKj$ . Then  $nHj$  sends  $Hj$  and  $C2$  to  $nLi$ .  $nLi$  decrypts  $C2$  to recover  $Hj$ ,  $Li$  and  $GKj$  by using  $kLi Hj$ . If  $nLi$  fails to check  $Hj$ ,  $Li$ , it discards the message and reports  $nHj$  to the BS as an illegitimate cluster head.

3) Membership Validation: After discovering all the neighboring nodes  $nLi$  ( $1 \leq i \leq n$ ) in the  $j$ th cluster,  $nHj$  computes  $C4 = EK0Hj (Hj, Mj)$  and transmits  $C4$  and  $Hj$  to the BS.

##### D. Key Update:

1) Pairwise Key Update: To update a pairwise encryption key, two nodes which shared the pairwise key perform a Pairwise Encryption Key Establishment process. On the other hand,

the pairwise master key does not require periodical updates, because it is not directly used to encrypt each session message.

2) Cluster Key Update: Only cluster head H-sensors can update their cluster key. If a L-sensor attempts to change the cluster key, the node is considered a malicious node.

#### E. Node Movement:

1) Node Leave: A node may leave a cluster due to node failure, location change or intermittent communication failure. There are both proactive and reactive ways for the cluster head to detect when a node leaves the cluster.

2) Node Join: Once the moving node  $n_{Lm}$  leaves a cluster, it may join other clusters or return to the previous cluster after some period. For the sake of simplicity, we assume that  $n_{Lm}$  wants to join the  $l$ th cluster or return to the  $j$ th cluster.

#### F. Key Revocation:

1) Compromised Node: The BS generates a CompNode message and a  $EK_{0H}(CompNode, L_c)$ . Then it sends  $EK_{0j}(CompNode, L_c)$  to all  $n_{Hj}$ , ( $1 \leq j \leq N_2$ ). After all cH-sensors decrypt the message, they update the revocation list of their clusters. Then, if related keys with  $n_{Lc}$  exist, the c related keys are discarded. Other than  $n_{Lc}$ ,  $n_{Hj}$  performs the cNode leave operations to change the current cluster key with c the remaining member nodes.

2) Compromised Cluster Head: After the BS generates a CompHeader message and a  $EK_{0Li}(CompHeader, H_j)$ , it sends the message to all  $n_{Li}$  ( $1 \leq i \leq n$ ) in the  $j$ th cluster. The BS also computes  $EK_{0Hi}(CompHeader, H_j)$ , ( $1 \leq i \leq N_2$ ,  $i \neq j$ ) and transmits it to all H-sensors except  $n_{Hj}$ .

#### G. Addition of a New Node:

Before adding a new node into an existing networks, the BS must ensure that the node is not compromised. Then new node  $n_{Ln+1}$  establishes a full private/public key through the node registration phase. Then, the public system parameters, a full private/public key and individual key  $K_{0Ln+1}$  are stored into  $n_{Ln+1}$ . The BS generates  $EK_{0Hj}(NewNode, L_{n+1}, pk_{Ln+1})$  and sends it to all  $n_{Hj}$ , ( $1 \leq j \leq N_2$ ).

## V. CONCLUSIONS

This paper propose the first certificateless effective key management protocol (CL-EKM) for secure communication in dynamic WSNs. CL-EKM supports efficient communication for key updates and management when a node leaves or joins a cluster and hence ensures forward and backward key secrecy. Our scheme is resilient against node compromise, cloning and impersonation attacks and protects the data confidentiality and integrity.

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